

Topology Control and Mobility Management in Mobile Ad Hoc Networks

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Mobile Ad Hoc Networks

- Infrastructure-less multi-hop wireless networks formed by mobile nodes
- Challenges
 - Frequent topology changes
 - Lack of central control
 - Limited network resources
 - Energy
 - Bandwidth
 - Computing power



Topology Control (TC)



- Goal

- Reduce transmission range (r) while maintaining network connectivity

- Motivation

- Energy efficiency

- $E_{Tx} \approx r^\alpha \quad (2 \leq \alpha \leq 4)$

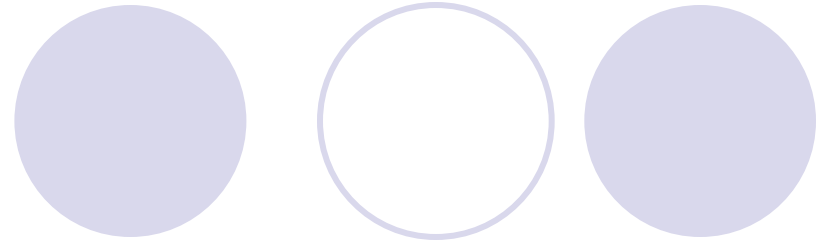
- $r=10$, single hop: $E_{Tx} \approx 1000 \quad (\alpha=3)$

- $r=1$, 10 hops: $E_{Tx} \approx 10$

- Channel spatial reuse

- Interference Area $\approx r^2$

TC (Continued)



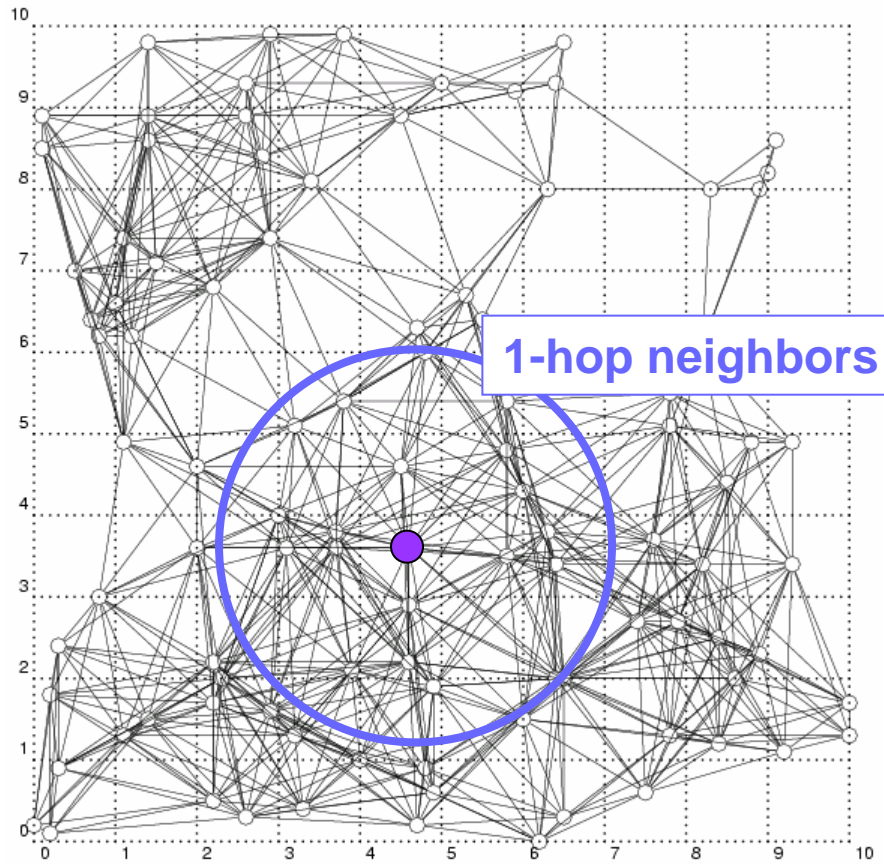
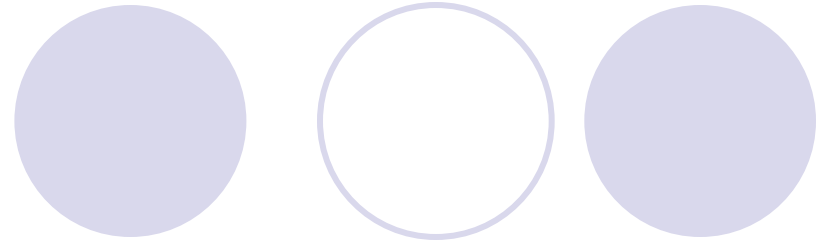
● Method

- Each node collects information of **1-hop neighbors** using “Hello” beacons
- Select a few **logical neighbors** from 1-hop neighbors
- Adjust transmission power to cover only logical neighbors

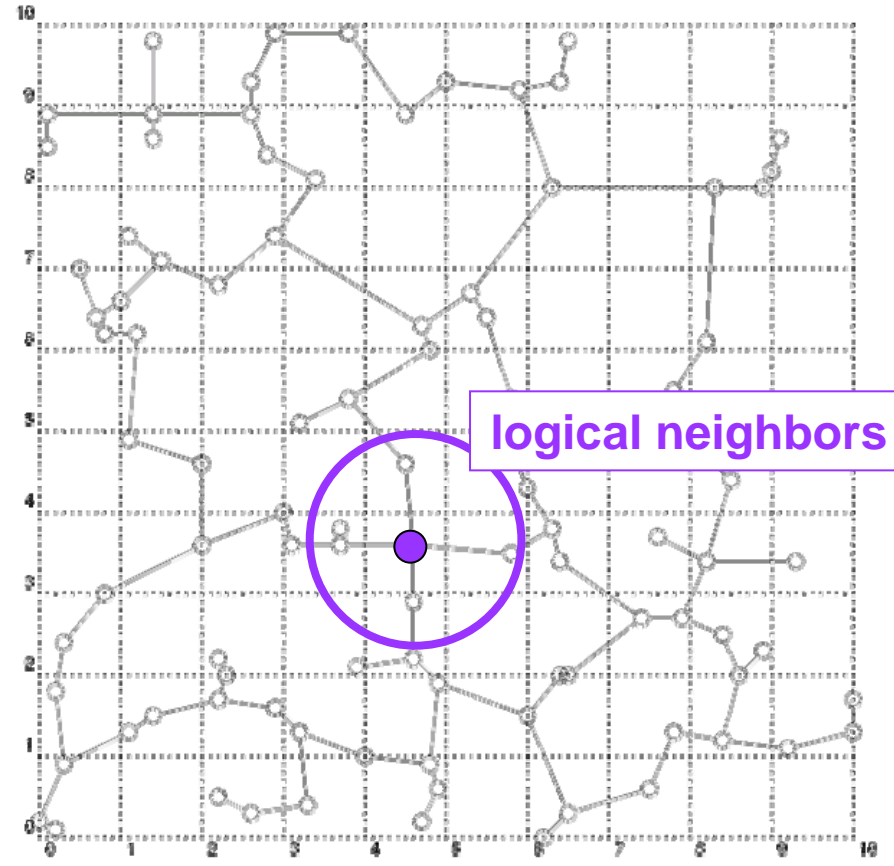
● Logical Topology

- Virtual network formed by **logical links** (i.e., links between logical neighbors)

TC Example



Original topology



Logical topology

Existing TC Protocols



- Select logical neighbors using local info.
 - Relative neighborhood graph (RNG)
 - Minimal energy mobile wireless network and extensions (LSPT)
 - Local minimum spanning tree (LMST)
 - Yao graph
 - Cone-based topology control (CBTC)

Relative Neighbor Graph (Toussaint, 1980)

- Remove link (u,v) if ...

- $\exists w$ s.t.

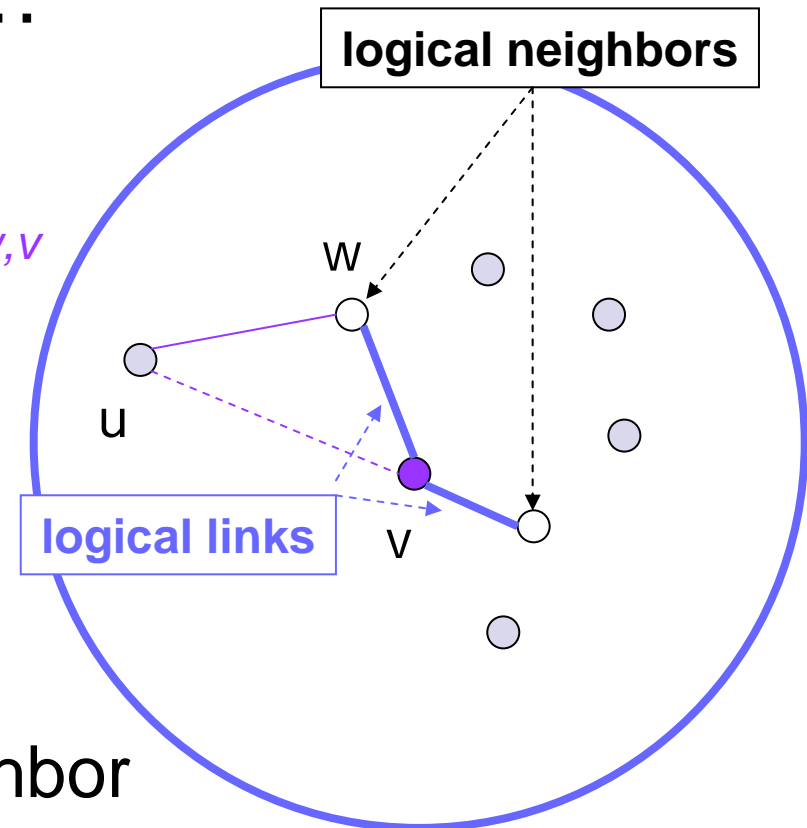
$$d_{u,v} > d_{u,w} \text{ and } d_{u,v} > d_{w,v}$$

- where

$$d_{u,v} = \text{dist. from } u \text{ to } v$$

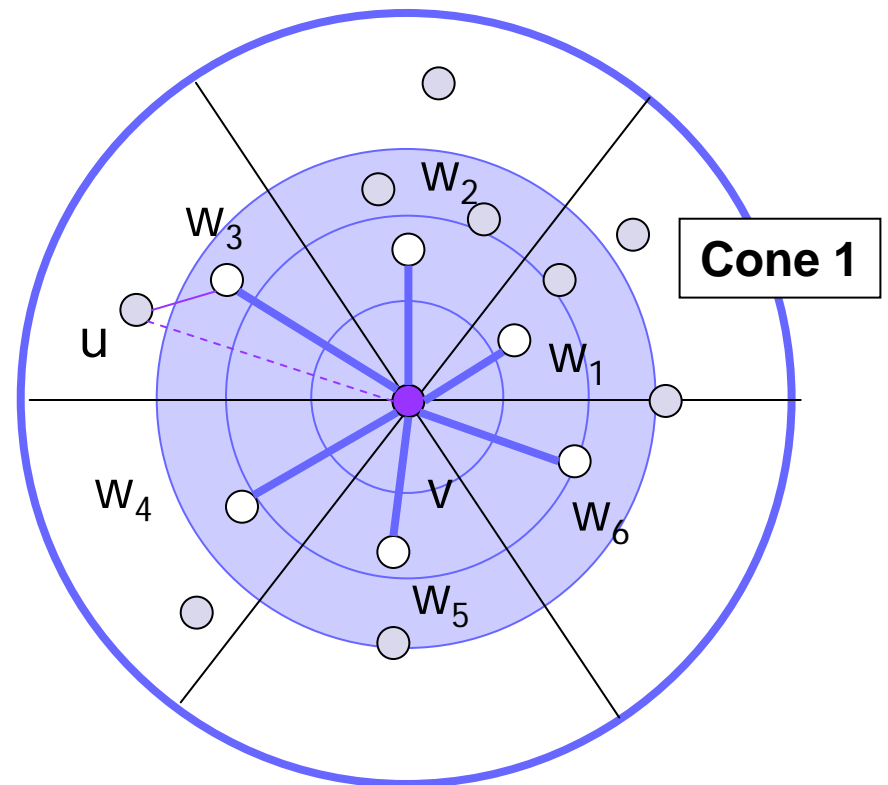
- From v 's view

- u is a non-logical neighbor



Yao Graph

- Divide neighborhood into 6 cones
- Select a nearest node from each cone as logical neighbor
- A special case of RNG
- Incremental search
 - Until one neighbor detected in each cone



Minimal Energy Mobile Wireless Network (Rodoplu and Meng, 1999)

- MEMWN

- Logical neighbors forms a closure

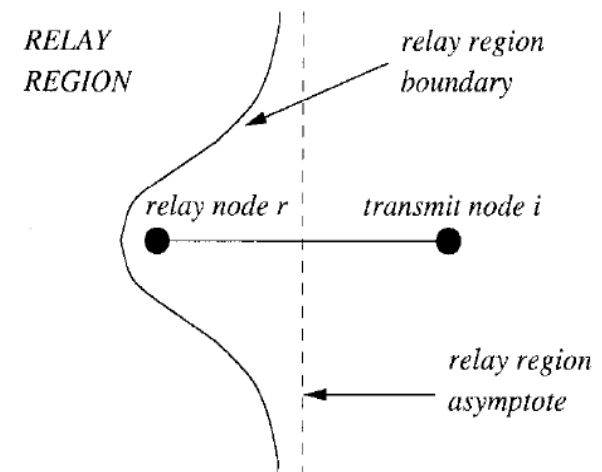
- \forall outsider j , \exists insider k

s.t. $E_{i,j} > E_{i,k} + E_{k,j}$

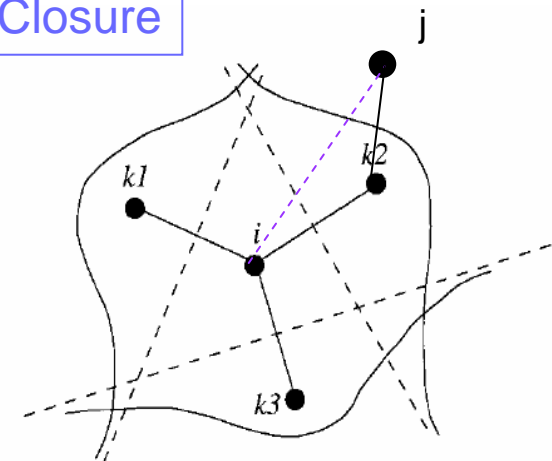
- Extension (Li and Halpern, 2001)

- Construct a shortest path from current node to each 1-hop neighbor w.r.t. Tx power

- First hop relays are logical neighbors

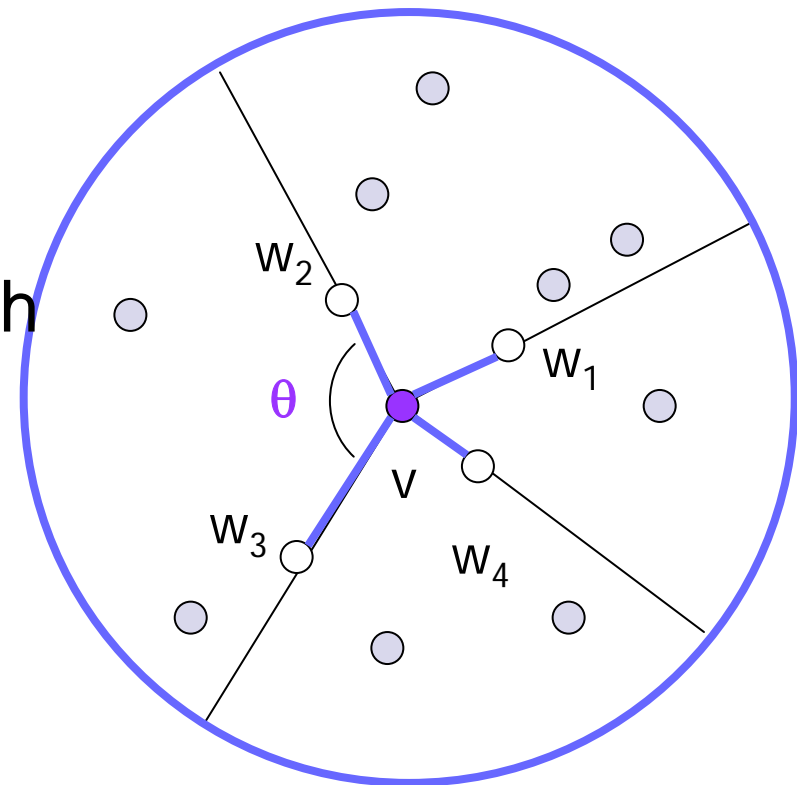


Closure



Cone-Based Topology Control (Li and Halpern, 2001)

- Find k logical neighbors that are
 - Nearest to the current node
 - The maximal cone width $\theta \leq 2\pi/3$
- Improve Yao graph
- Extension
 - $\theta \leq 4\pi/5$ with some additional constraint



Local Minimal Spanning Tree (Li, Hou, and Sha, 2003)

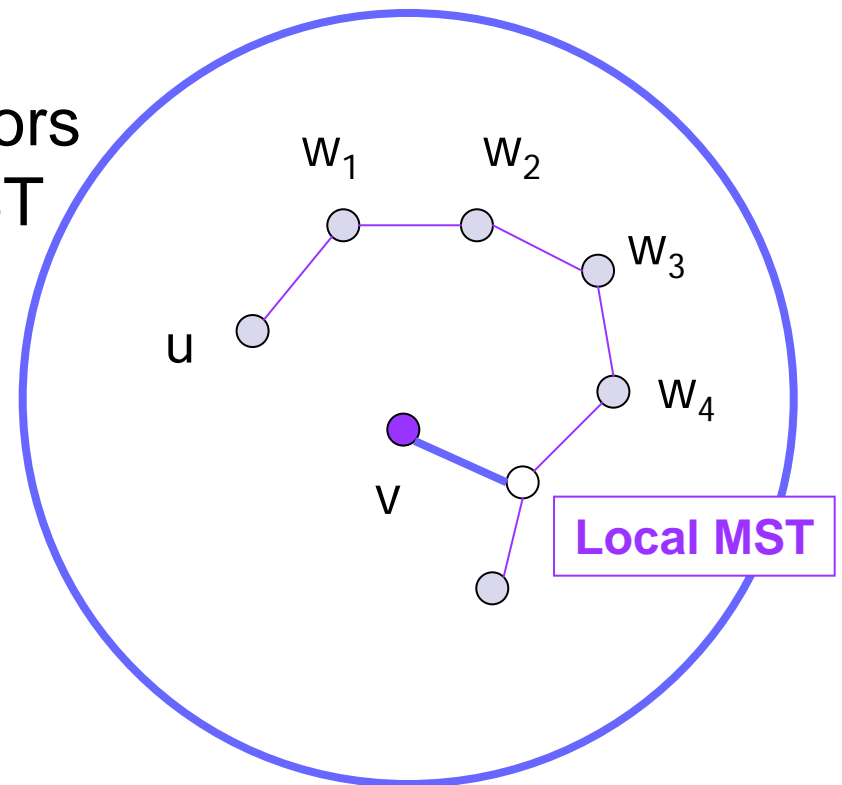
- LMST

- Formed by shortest links
- Connect all 1-hop neighbors
- Logical neighbors are MST neighbors

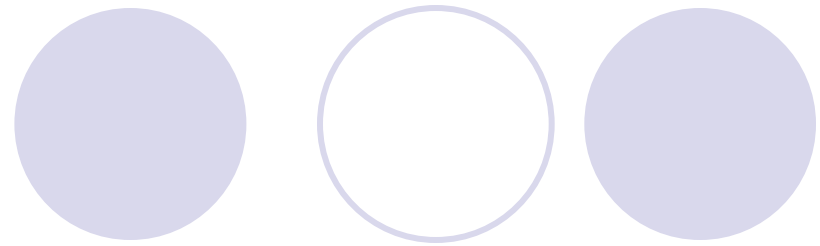
- In other words ...

- Remove link (u,v) if
exist w_1, w_2, \dots, w_k s.t.

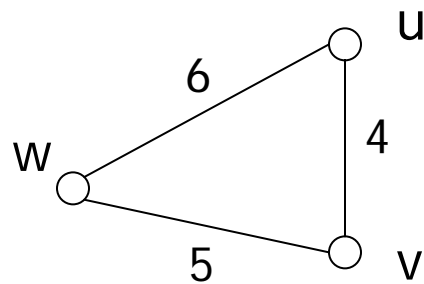
- $d_{u,v} > d_{u,w_1}$
- $d_{u,v} > d_{w_1,w_2}$
- ...
- $d_{u,v} > d_{w_k,v}$



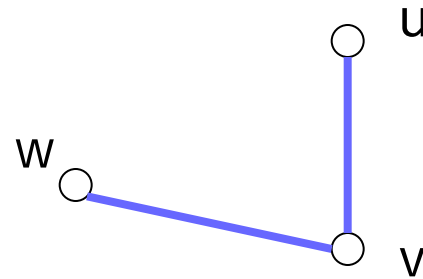
LMST Example



Original topology

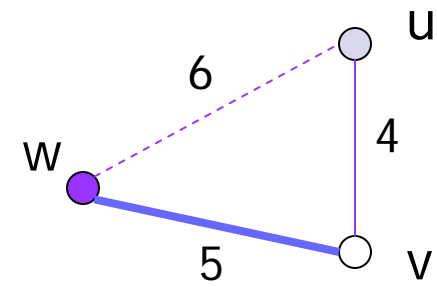
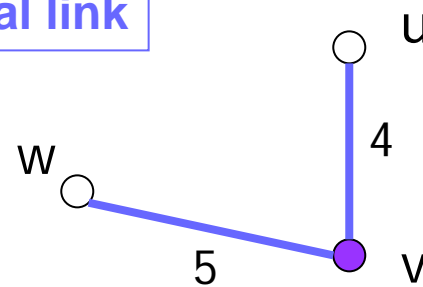
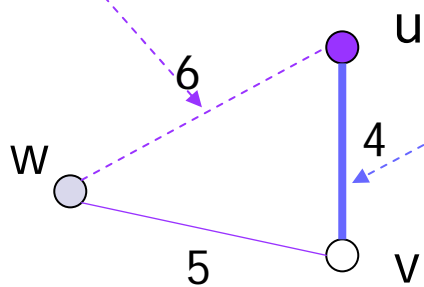


Logical topology



removed link

logical link



Local MSTs at nodes u , v , and w

Summary: Link Removal Rule

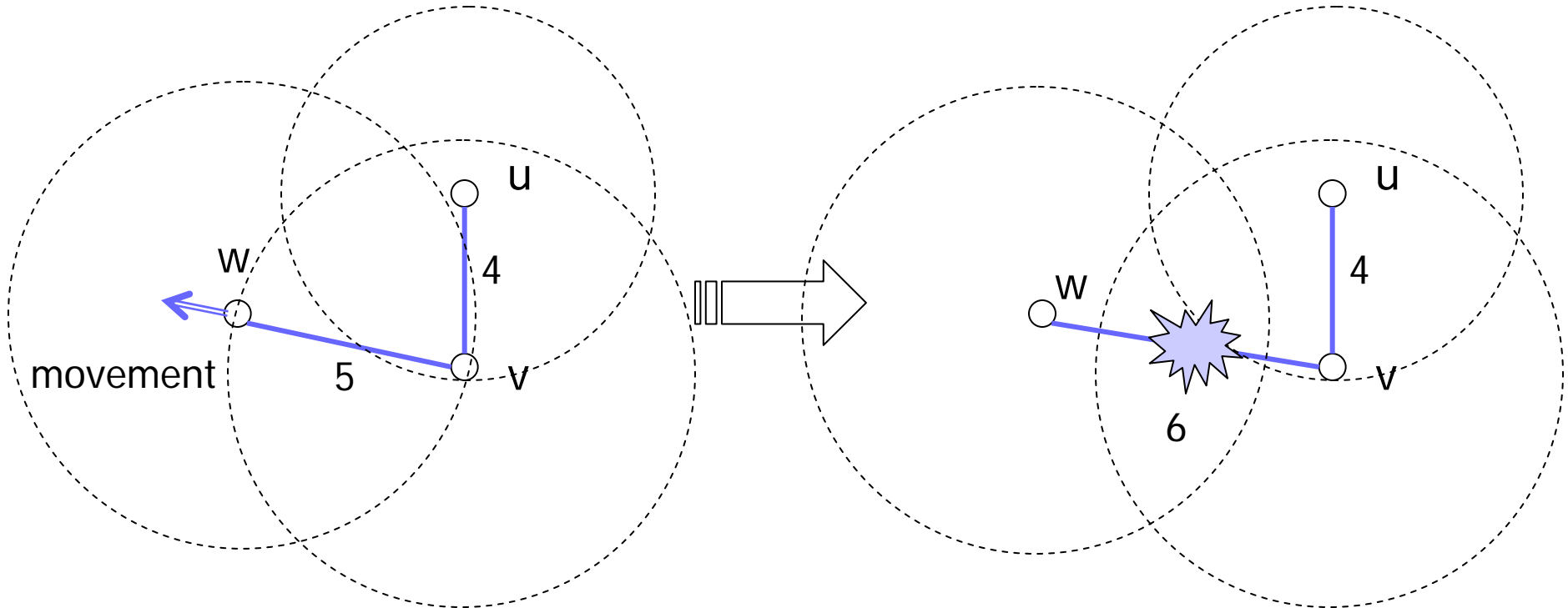
- Detect loops $(u, w_1, w_2, \dots, w_k, v)$ within 1-hop neighborhood
- Break loop by removing a link (u, v) with the highest cost $C_{u,v}$
- Some criteria break more loops than the others
- Guarantees connectivity in static networks
- What if the network is not static?

TC in Mobile Networks



- Inaccurate position information
 - Actual transmission range may not cover all logical neighbors
 - Solution: slightly increase the actual transmission range (**buffer zone**)
- Inconsistent local view
 - Simultaneous removal => disconnection
 - Solution: enforce **view consistency**

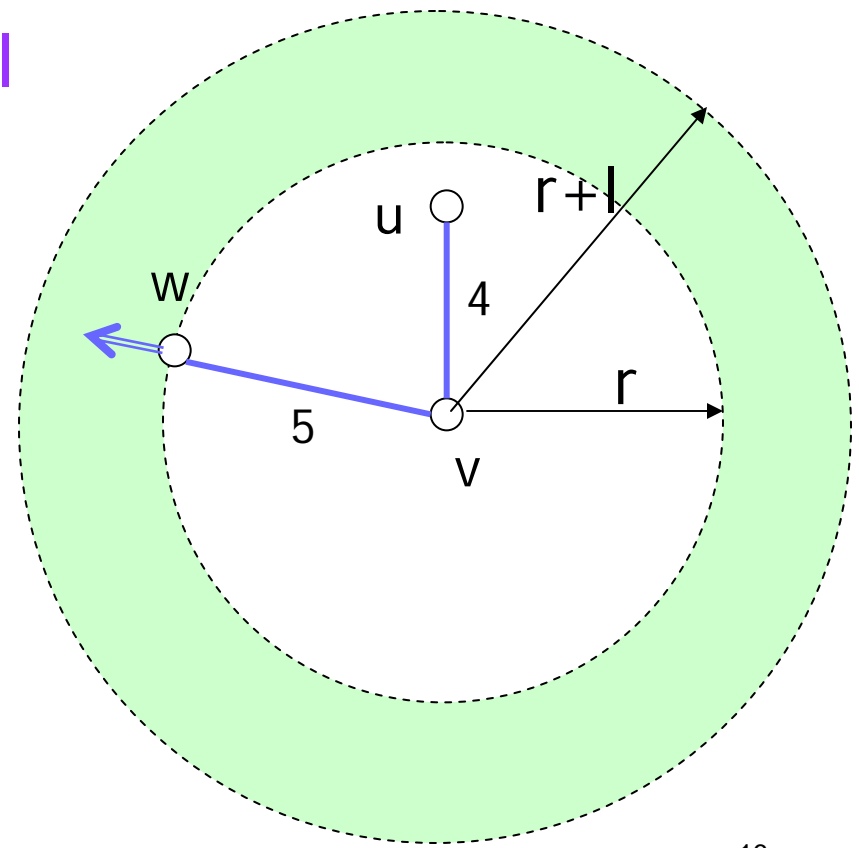
Outdated Information



- Link (v,w) is broken after w moves out of v 's transmission range

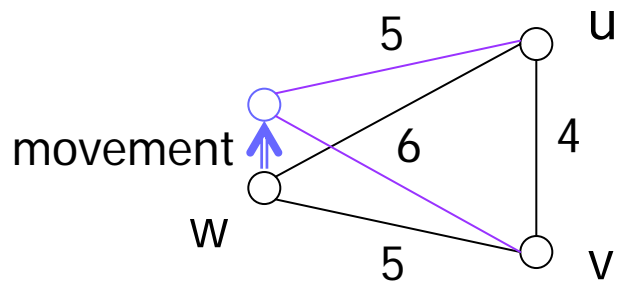
Solution: Buffer Zone

- Increase the actual transmission range r by l to create a buffer zone that tolerate node movement
- $l = dt$, where t is the maximal relative moving speed, and d is the maximal delay
- Using a $l < dt$ is possible

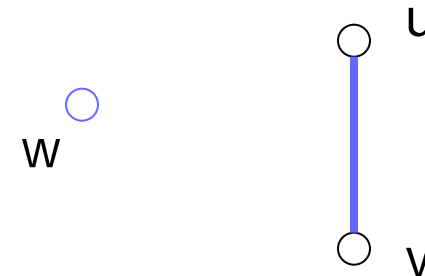


Inconsistent Local Views

Original topology

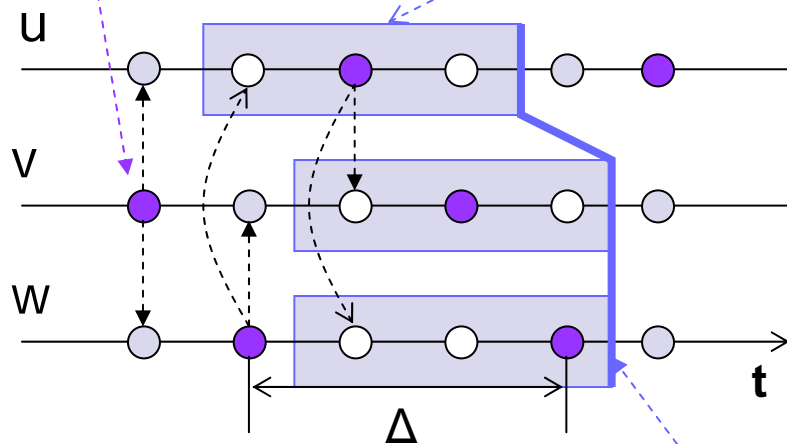


Logical topology



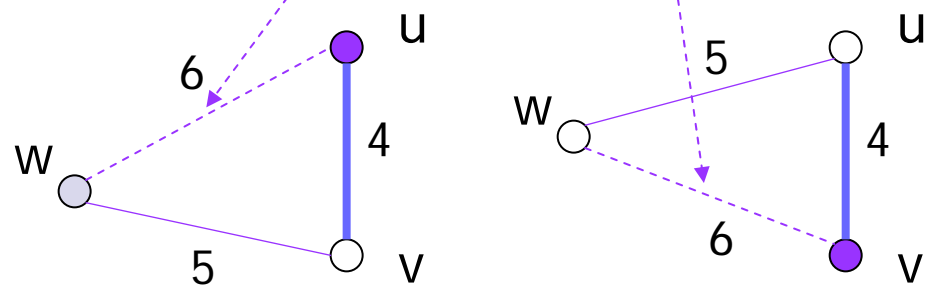
"Hello" msg

Local view



Decision making times

simultaneous removal



Local MSTs at nodes u and v

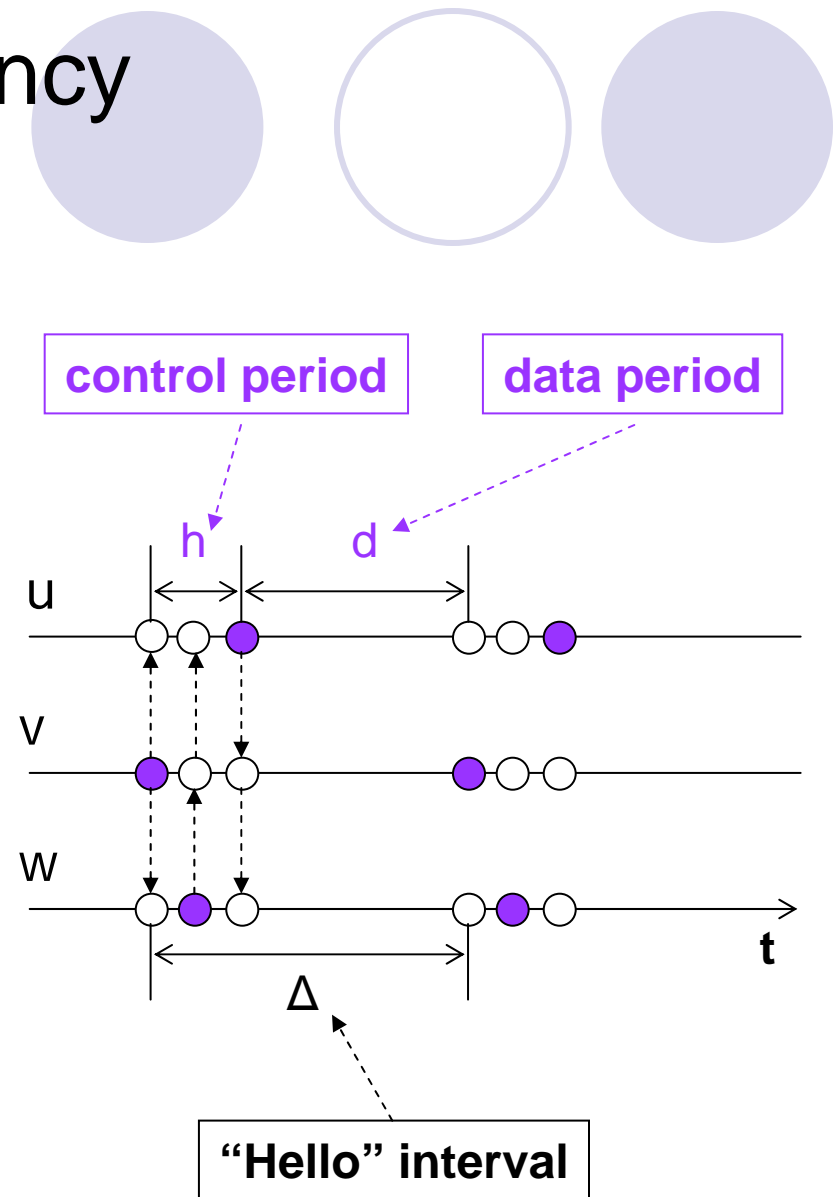
Strong View Consistency (Wu and Dai, 2004)

- Solution

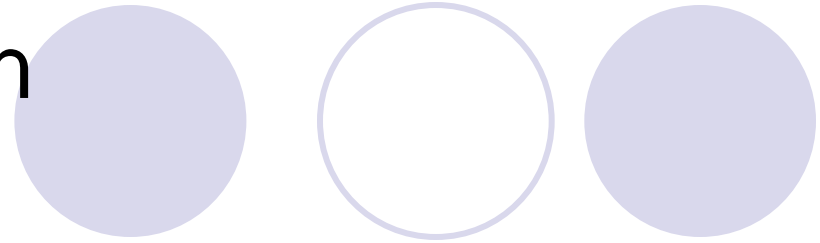
- All nodes exchange “hello” messages at control period (h)
- Decisions are made at the beginning of each data period (d)

- Problems

- “Hello” message collision
- Require synchronous clocks



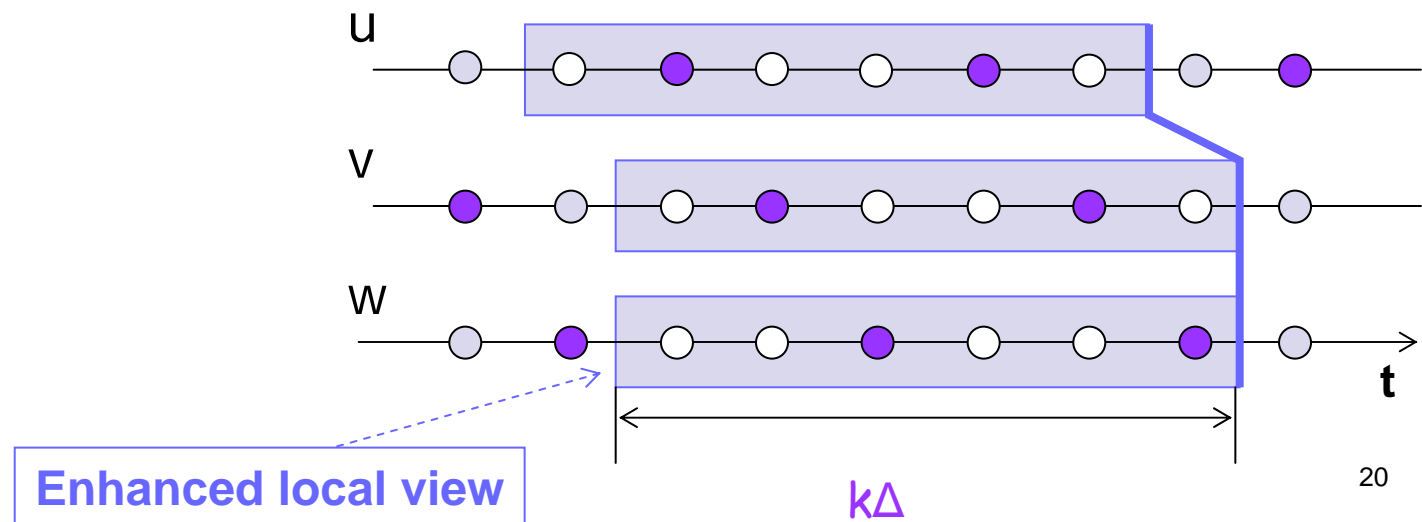
Asynchronous Solution (Dai and Wu, 2005)



- Enhanced local view
 - Use history information
- Enhanced link removal rule
 - Make conservative decisions
- Weak view consistency
 - Guarantee connectivity in spite of mobility and asynchronous decision making

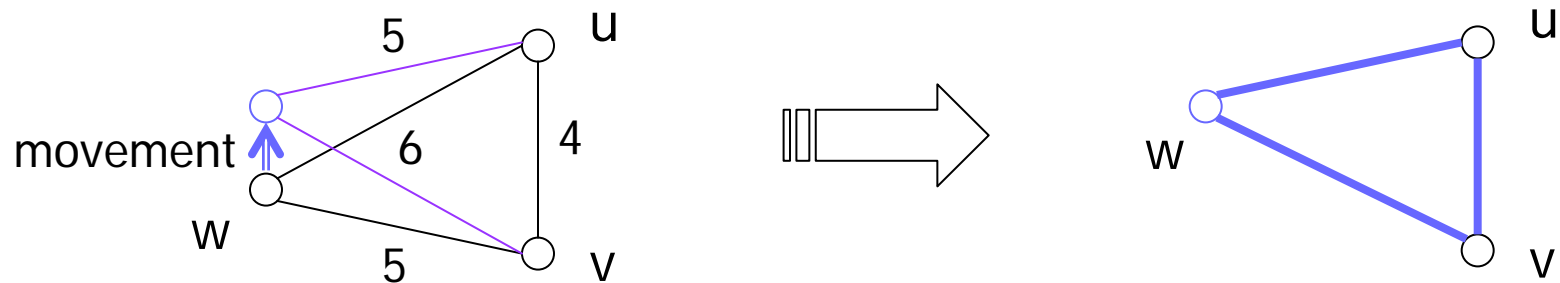
Enhanced Local View

- Contain k recent "Hello" messages
- Up to k^2 cost values for each link (u,v)
 - $Max(C_{u,v})$: maximum cost in current local view
 - $Min(C_{u,v})$: minimum cost in current local view
 - $MinMax(C_{u,v})$: minimum $Max(C_{u,v})$ in all local views
 - $MaxMin(C_{u,v})$: maximum $Min(C_{u,v})$ in all local views



Enhanced Link Removal Rule

- A link (u,v) can be removed, if
 - a loop $(u, w_1, w_2, \dots, w_k, v)$ exist, and
 - $Min(C_{u,v}) > Max(C_{u,w_1}, Max(C_{w_1,w_2}), \dots, Max(C_{w_k,v}))$
- Previous example reconsidered
 - Cannot remove (u,w) : $Min(C_{u,w}) < Max(C_{v,w})$
 - Cannot remove (v,w) : $Min(C_{v,w}) < Max(C_{u,w})$



Weak View Consistency

- Local views are **weakly consistent**, if
 - $MinMax(C_{u,v}) \geq MaxMin(C_{u,v})$ for all (u,v)

Theorem: Applying enhanced link removal rule using weakly consistent local views guarantees connectivity

Theorem: Enhanced local views containing two recent “Hello” messages are weakly consistent.

- Put together
 - Two “Hello” messages from each neighbor are enough to guarantees connectivity

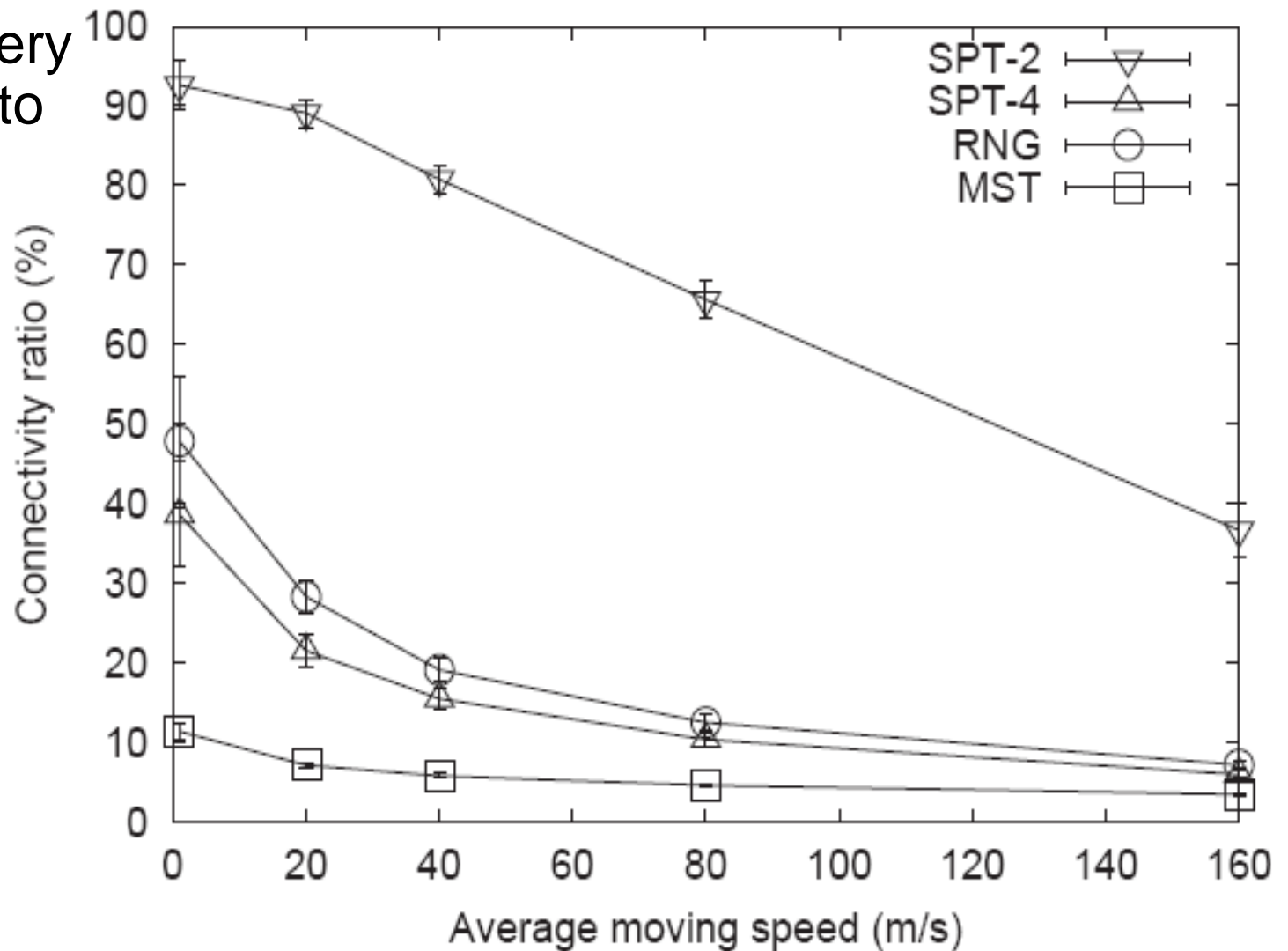
Simulation

A decorative graphic at the top of the slide consists of two groups of circles. The first group on the left has a solid light purple circle on the left and an outlined light purple circle on the right. The second group on the right has a solid light purple circle on the left, an outlined light purple circle in the middle, and a solid light purple circle on the right.

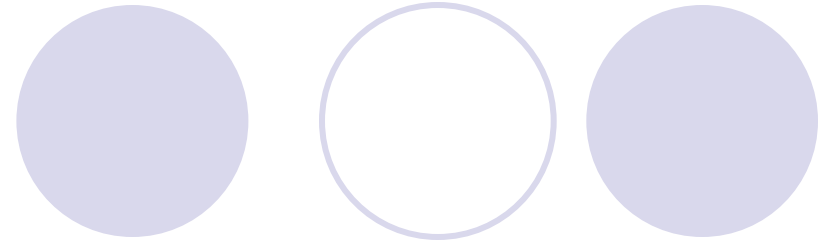
- Simulator: ns2
- Simulated algorithms
 - MST: local minimal spanning tree
 - RNG: relative neighborhood graph
 - SPT: local shortest spanning tree with link cost d^2 (a=2) and d^4 (a=4), where d is distance.
- Mobility model
 - Random waypoint, avg. speed 1-160m/s
- Ideal MAC layer without collision/contention
- Connectivity ratio:
(pairs of connected nodes) / (total node pairs)

TC Protocols Under Mobile Environment

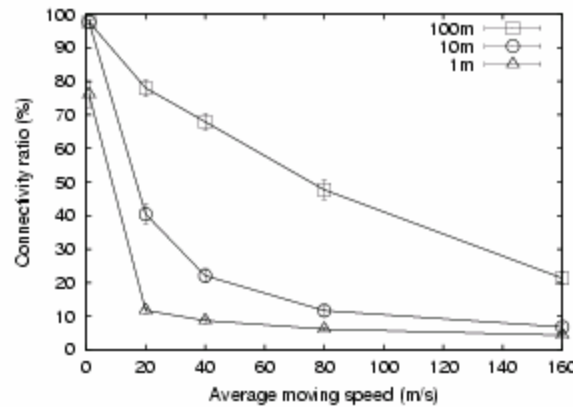
- Fragile: very sensitive to mobility



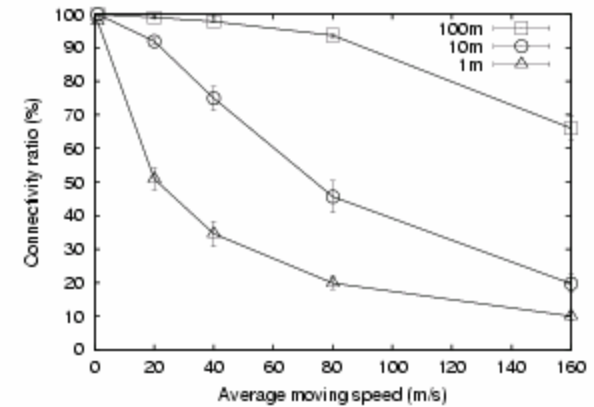
Buffer Zone Width



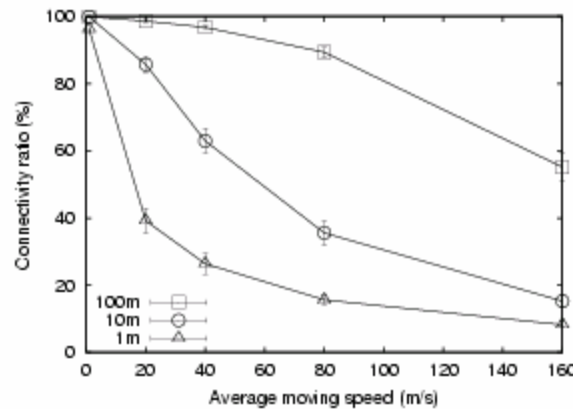
- Connectivity ratio increased as buffer zone width increases
- Buffer zone alone does not guarantee connectivity



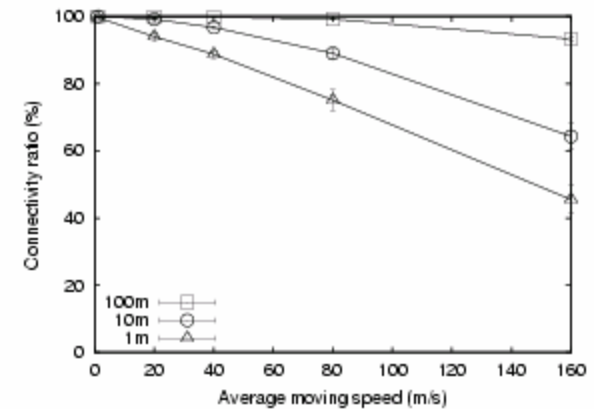
(a) MST



(b) RNG



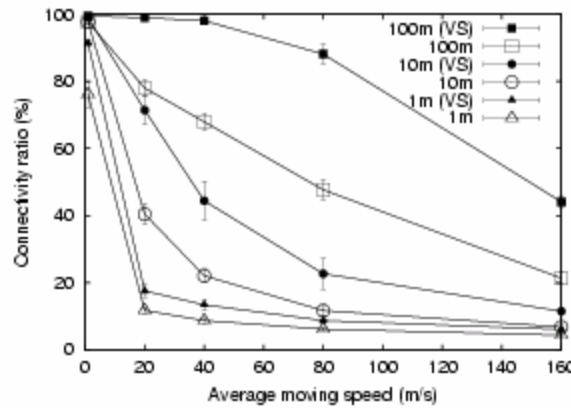
(c) SPT ($\alpha = 4$)



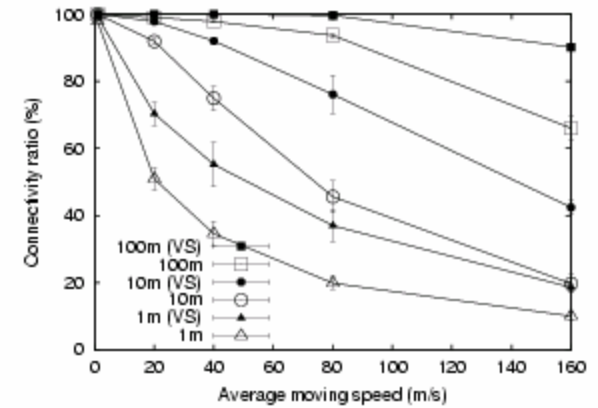
(d) SPT ($\alpha = 2$)

Consistent Views

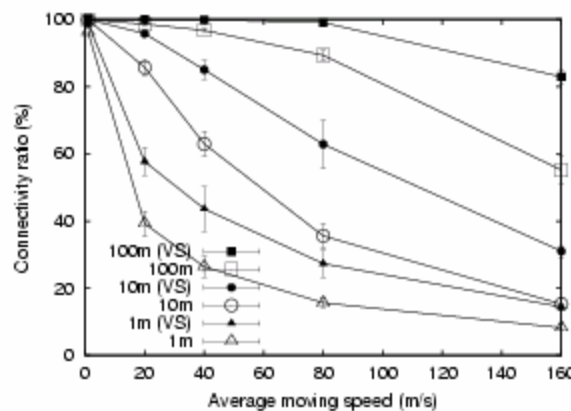
- Using consistent views improves connectivity ratio significantly



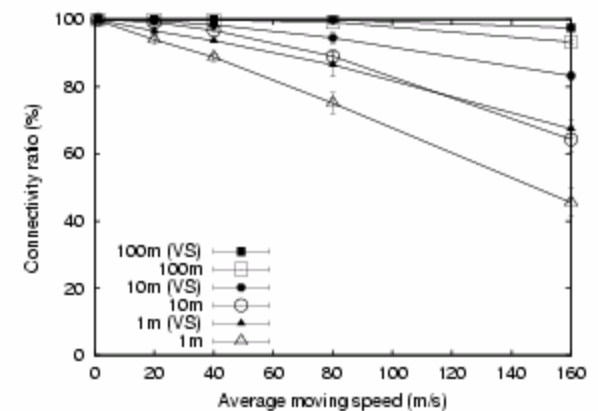
(a) MST



(b) RNG



(c) SPT ($\alpha = 4$)



(d) SPT ($\alpha = 2$)

Conclusion



- In MANETs, TC may cause
 - Insufficient actual transmission power
 - Disconnected logical topology
- Weak consistency scheme
 - Guarantee connectivity using 2 recent “Hello” messages
 - No synchronization overhead
 - Slightly increase the number of logical links
 - Enhance many existing TC protocols
- Future directions
 - Fault tolerance against collision