

Energy-Efficient Continuous and Event-Driven Monitoring

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Outline

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Maximum Sensor Network Lifetime Problem

A formal definition of the energy preserving scheduling problem

- ✓ **Sensor cover** : A set of sensors C covering R .
- ✓ **A monitoring schedule**: a set of pairs $(C_1, t_1), \dots, (C_k, t_k)$.
 - C_i is a sensor cover;
 - t_i is time during which C_i is active.

Maximum Sensor Network Lifetime problem

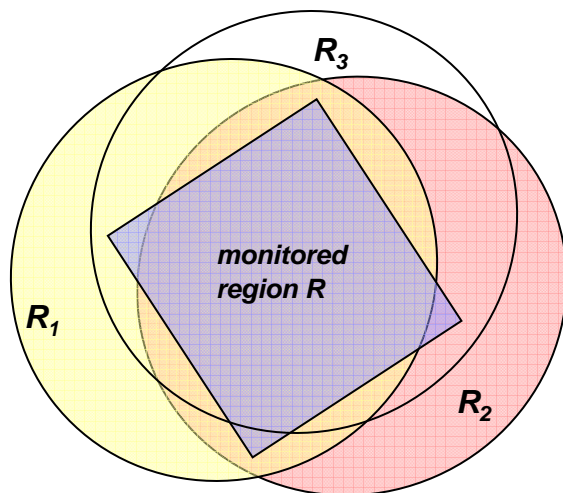
Given: a monitored region R , a set of sensors p_1, \dots, p_n , and monitored region R_i , and energy supply b_i for each sensor

Find: a *monitoring schedule* $(C_1, t_1), \dots, (C_k, t_k)$ with the *maximum length* $t_1 + \dots + t_k$, such that for any sensor p_i the total active time does not exceed b_i .

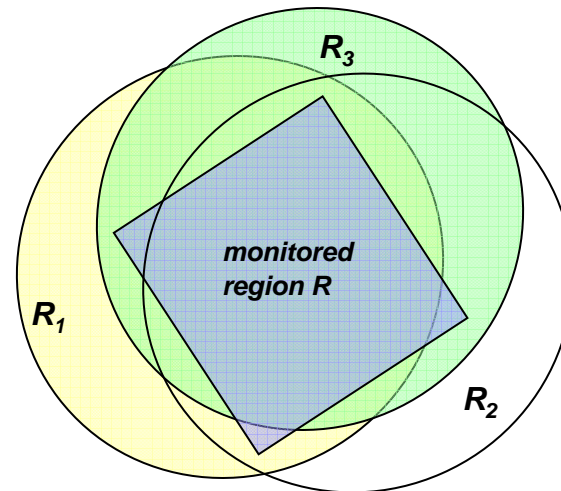


Example of Maximum Sensor Network Lifetime Problem

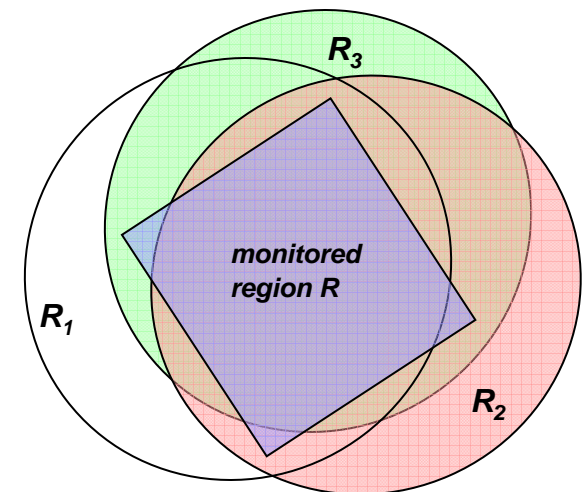
Advantage of switching between sensor covers:



sensors p_1 and p_2
for 1 time unit



sensors p_1 and p_3
for 1 time unit



sensors p_2 and p_3
for 1 time unit

Non-disjoint set covers:

- the schedule $(\{p_1, p_2\}, 1), (\{p_2, p_3\}, 1), (\{p_3, p_1\}, 1)$;
- **3 units of time.**



CONTINUOUS AND EVENT-DRIVEN SENSOR NETWORK MODEL

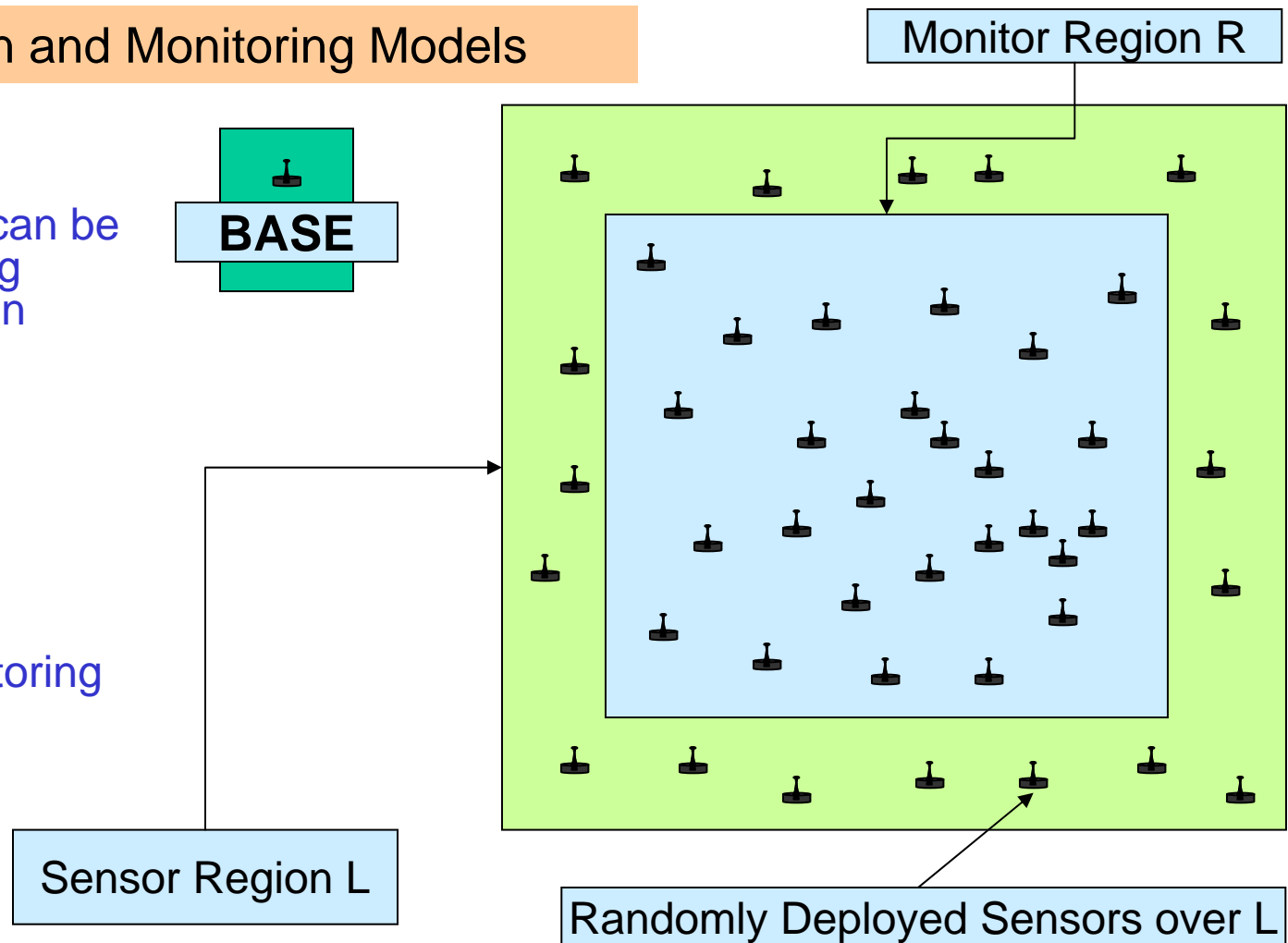
- **Given the regions which are required to monitor (or, in general, set of required targets)**
- **sensors who can monitor these targets**
- **energy supply**
- **energy consumption**
- **rate for monitoring**
- **listening and idle modes**
- **energy consumption for receiving and transmitting a package**
- **we explore the problem of maximizing sensor network lifetime**
- **sensors can interchange idle and active modes both for monitoring and communicating.**



CONTINUOUS AND EVENT-DRIVEN SENSOR NETWORK MODEL

Communication and Monitoring Models

- Each sensor can be in the following communication modes:
 - sleeping,
 - listening,
 - receiving and
 - sending.
- and two monitoring modes:
 - idle and
 - active.



- The set of sensors largely exceeds the necessary amount to monitor R

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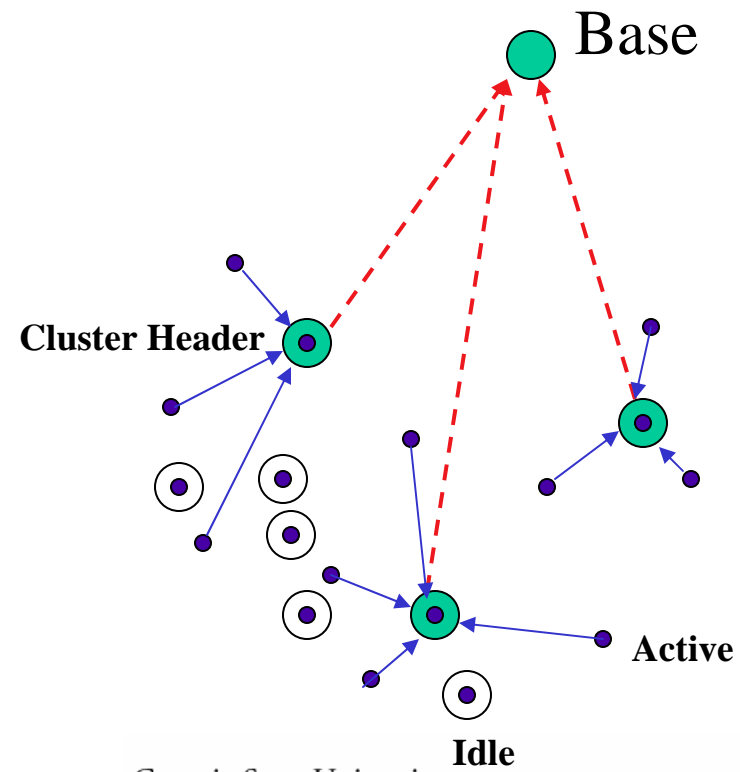
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(DEEPS) General overview

(DEEPS) Deterministic Energy-Efficient Protocol for Sensor networks target-monitoring protocol, system lifetime increase in **2 times!!!**

full-fledged simulation of the monitoring protocols on **NS2** combined with **LEACH** as a communication protocol



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Cluster-Based Communication

A Simple Algorithm

The problem: Select j cluster-heads of N nodes without communication among the nodes

The simplest solution:

- Each node determines a random number x between 0 and 1
- If $x < j / N \rightarrow$ node becomes cluster-head

...it's good, but:

Cluster-heads dissipate much more energy than non cluster-heads!

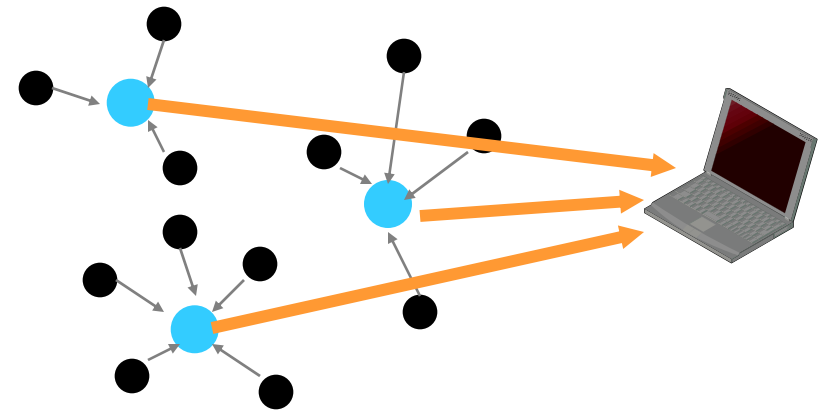
How to distribute energy consumption?



LEACH Communication Protocol

Low-Energy Adaptive Clustering Hierarchy

- Cluster-based communication protocol for sensor networks, developed at MIT
- Adaptive, self-configuring cluster formation



- The operation of LEACH is divided into rounds
- During each round a different set of nodes are cluster-heads
- Each node n determines a random number x between 0 and 1
- If $x < T(n) \rightarrow$ node becomes cluster-head for current round



(LBP) Load-Balancing Protocol for Sensing

(1) (*on-rule*)

whenever a sensor s has a target covered solely by itself (no alert- or on sensor covers it), s switches itself on, i.e., changes its state to "on".

(2) (*off-rule*)

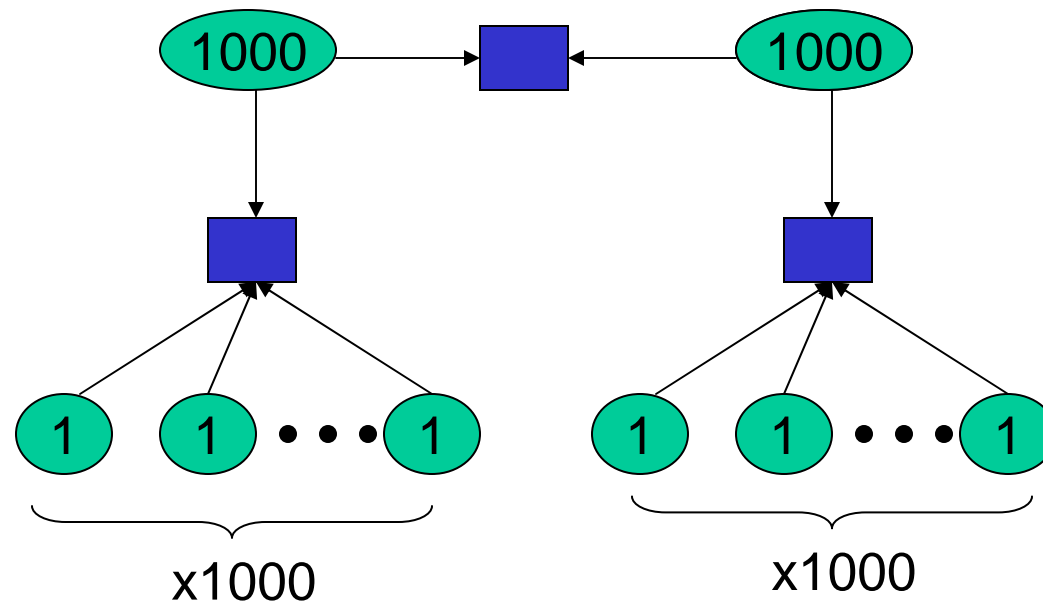
whenever each target covered by a sensor s is also covered either by an on sensor or an alert-sensor with a larger battery supply, s switches itself on, i.e., s changes its state to "off".

Each global reshuffle of LBP needs 2 broadcasts (to the neighbors) from each sensor and the resulted set of all active sensors form a minimal sensor cover.

The LBP is a reliable protocol.



(LBP) Bottleneck



- LBP time schedule is twice shorter since it uses the 1000-battery sensors simultaneously for 999 time units

Deterministic Energy-Efficient Protocol for Sensing

Algorithm

(1) (*on-rule*)

whenever a sensor s has a target covered solely by itself (no alert- or on-sensor covers it), s switches itself on, i.e., changes its state to "on".

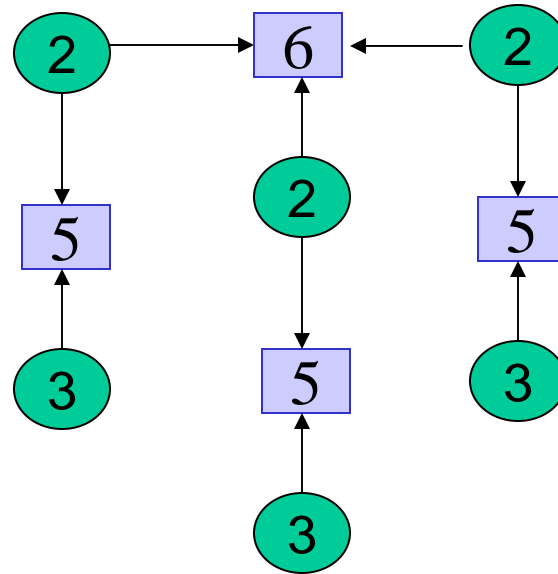
(2) (*off-rule*)

whenever a sensor s is not in charge of any target except those already covered by on-sensors, s switches itself off, i.e., changes its state to "off".

DEEPS is a reliable protocol. Each global reshuffle of DEEPS needs 2 broadcasts (to the 2-neighbors) from each sensor and the resulted set of all active sensors form a minimal sensor cover.



(DEEPS) Bottleneck



An example of reliability violation (circles are sensors and rectangles are targets, numbers correspond to battery supply). 3 lower sensors have 3 batteries each and the 3 upper sensors have 2 batteries each. Therefore, 3 lower targets are sinks with 5 batteries each. The upper target will be abandoned if all three upper poorer sensors will be switched off simultaneously.



NS2+LEACH Monitoring Simulations

Environment : NS2 – Network Simulator

LEACH communication protocol

DEEPS - Deterministic Energy-Efficient Protocol for Sensing

LBP – Load-Balancing Protocol for Sensing

1-DEEPS which is DEEPS but with a single reshuffle and local reparation on node die

EUPS - Energy Unaware Protocol for Sensing – where all sensors continuously monitor their targets



Simulation Results

Results which are represented in this presentation are obtained for 3 scenarios:

Scenario-1:

- Square territory $100m \times 100m$ which is divided into the small square faces $1m \times 1m$, and each face is considered like a target with coordinates equals with the middle of the face.
- 10.000 targets – good approximation for real area.
- Random distribution of sensors
- Faces are covered by one or more sensors, sensing radius is 5m.

Scenario-3:

- The same as Scenario-1, with additional restriction: All faces are covered by at least 3 sensors.

Scenario-2:

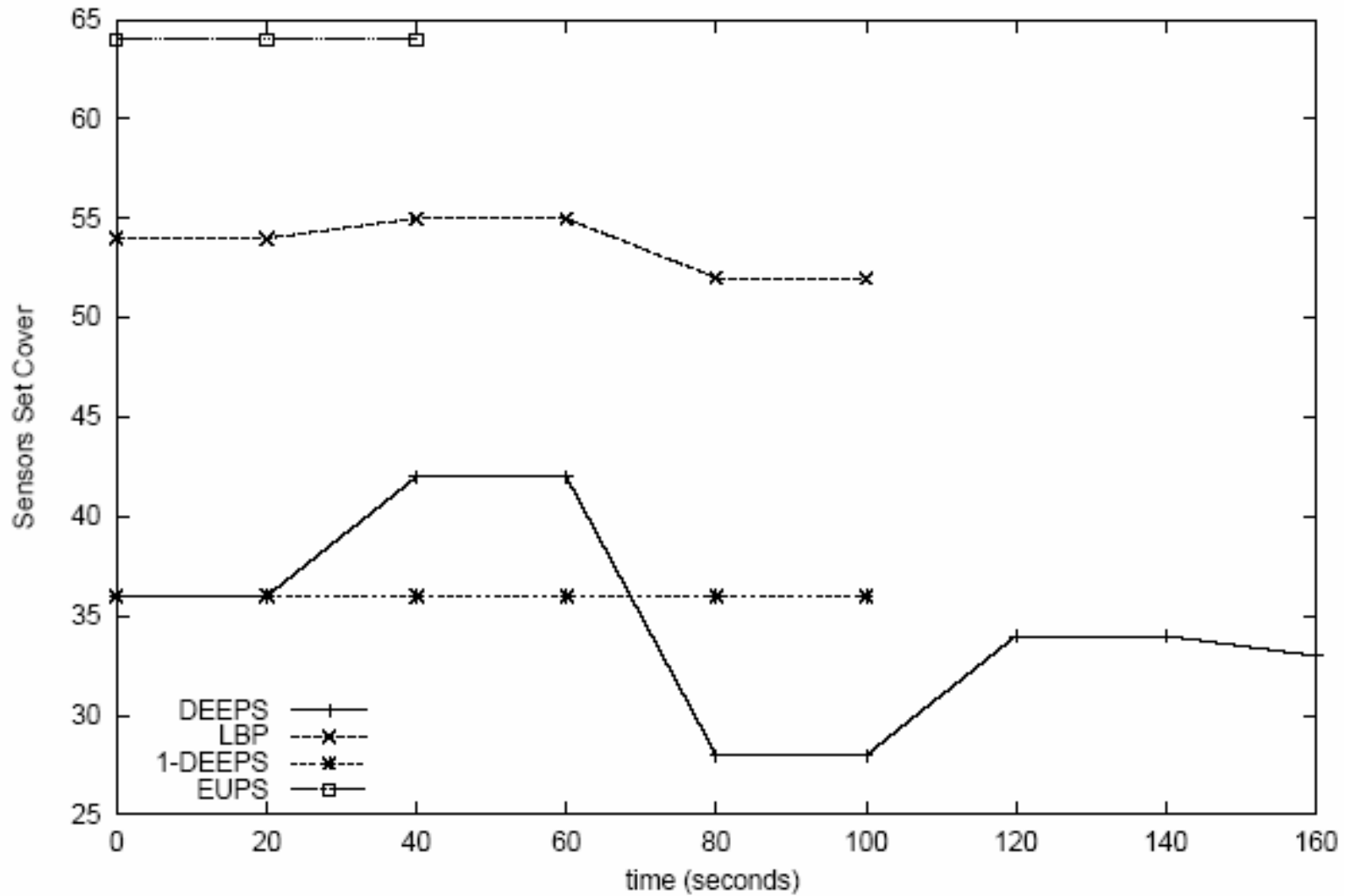
- Random distribution of 100.000 targets in $100m \times 200m$.
- All others are the same as in Scenario-3

Experimental results are for constant initial energy distribution $4(J)$ or random between 1 and $4(J)$.



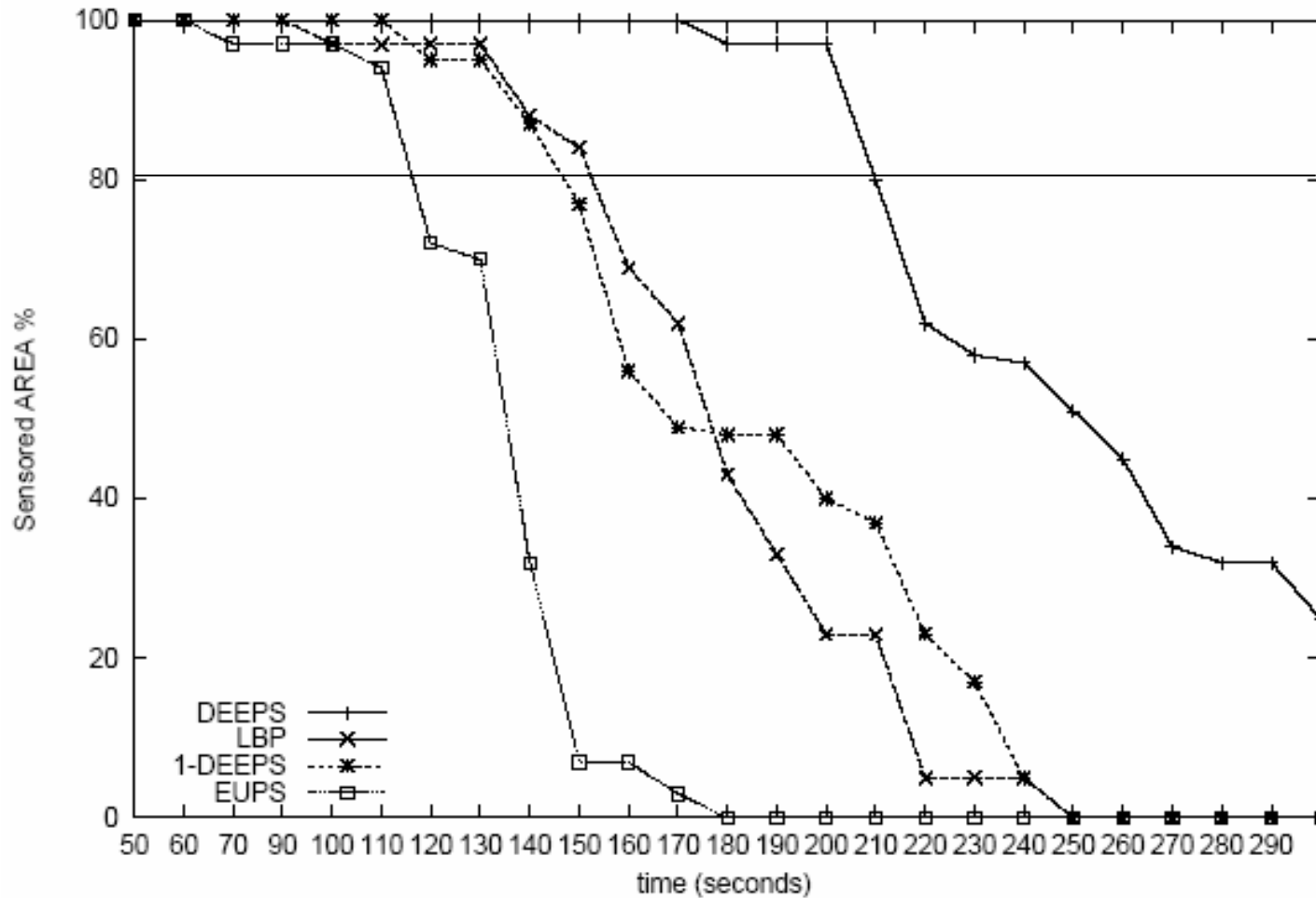
Scenario-1

Number of active sensors



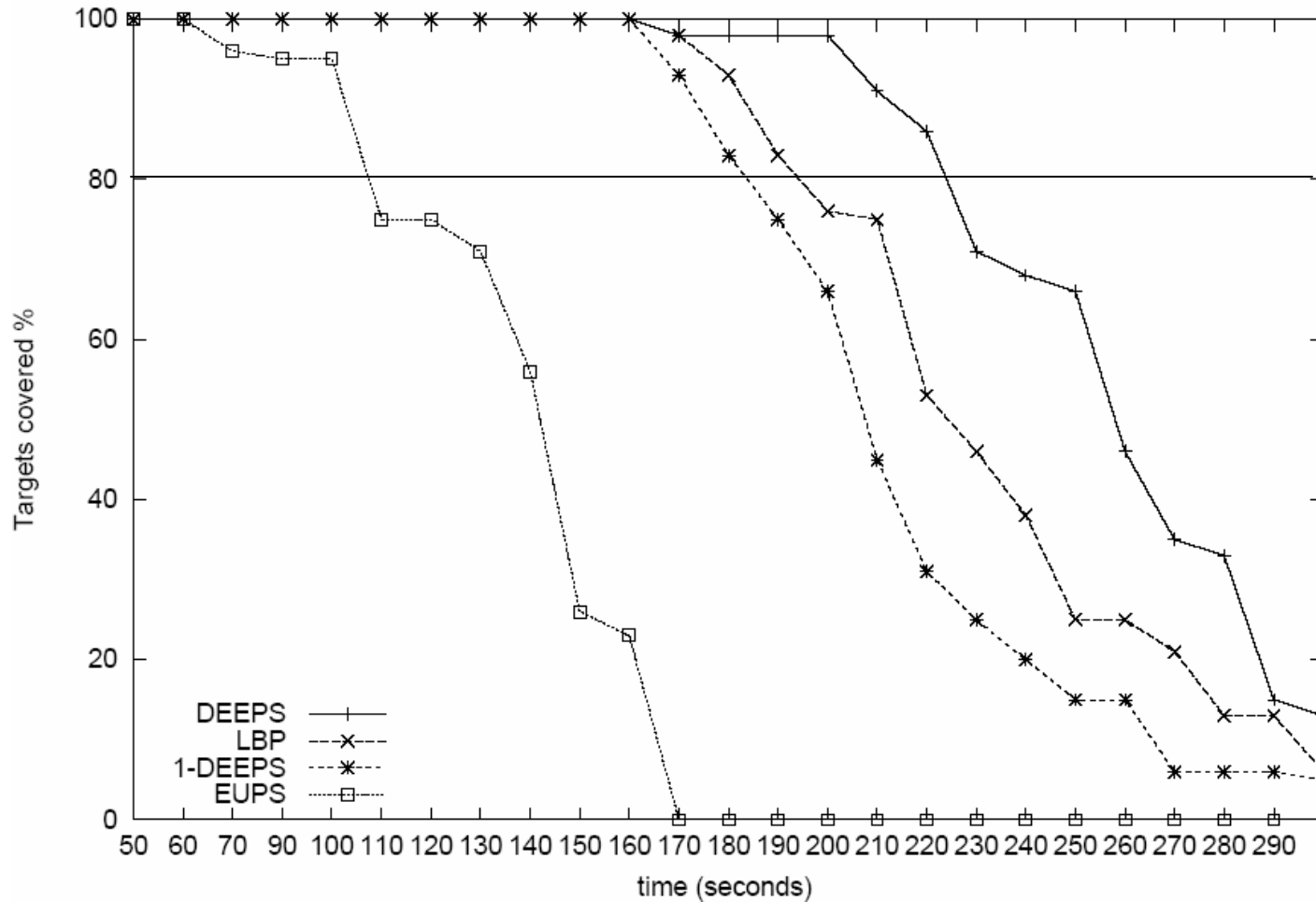
Scenario-1

Covered area %



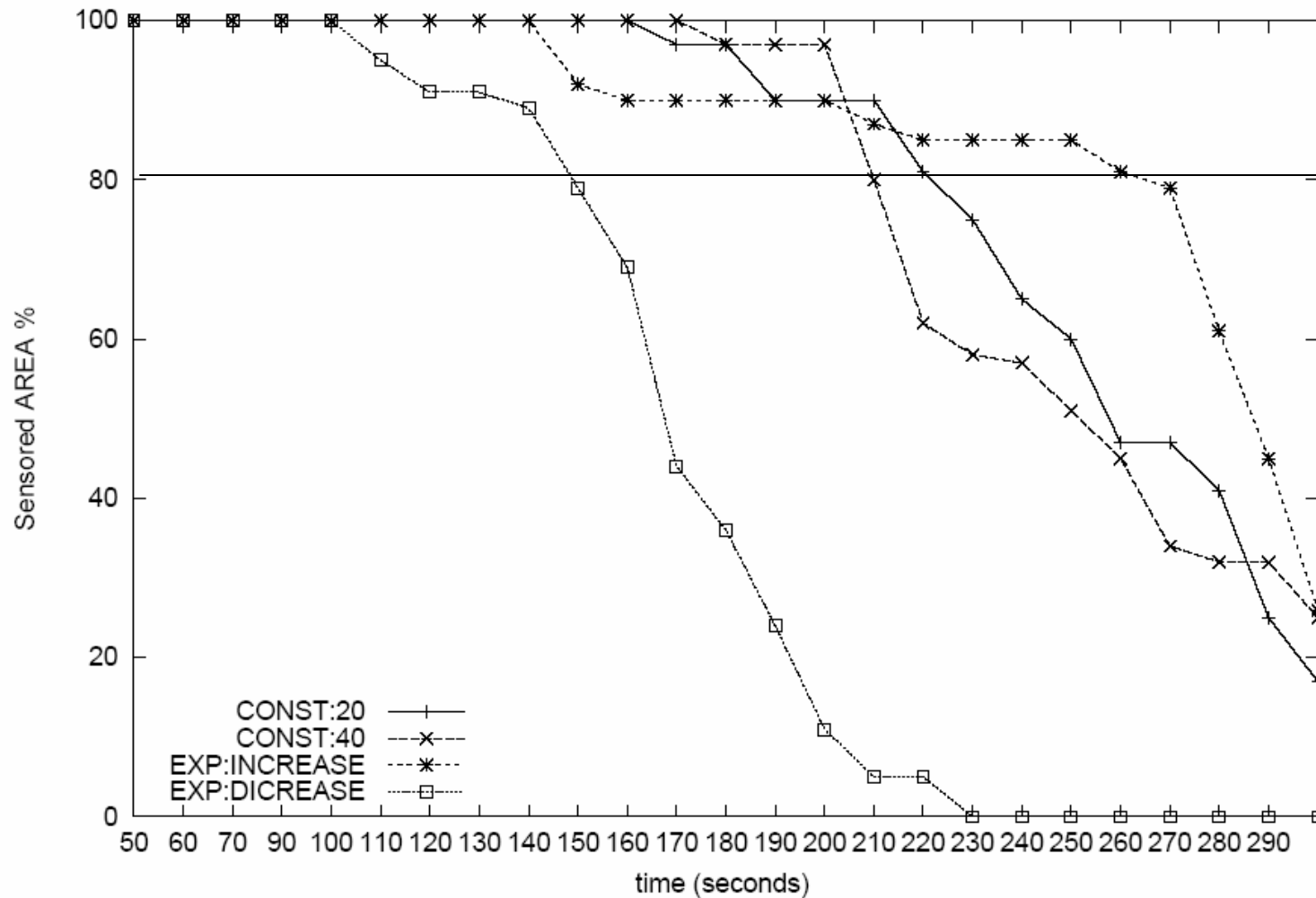
Scenario-3 TARGET

Targets covered %



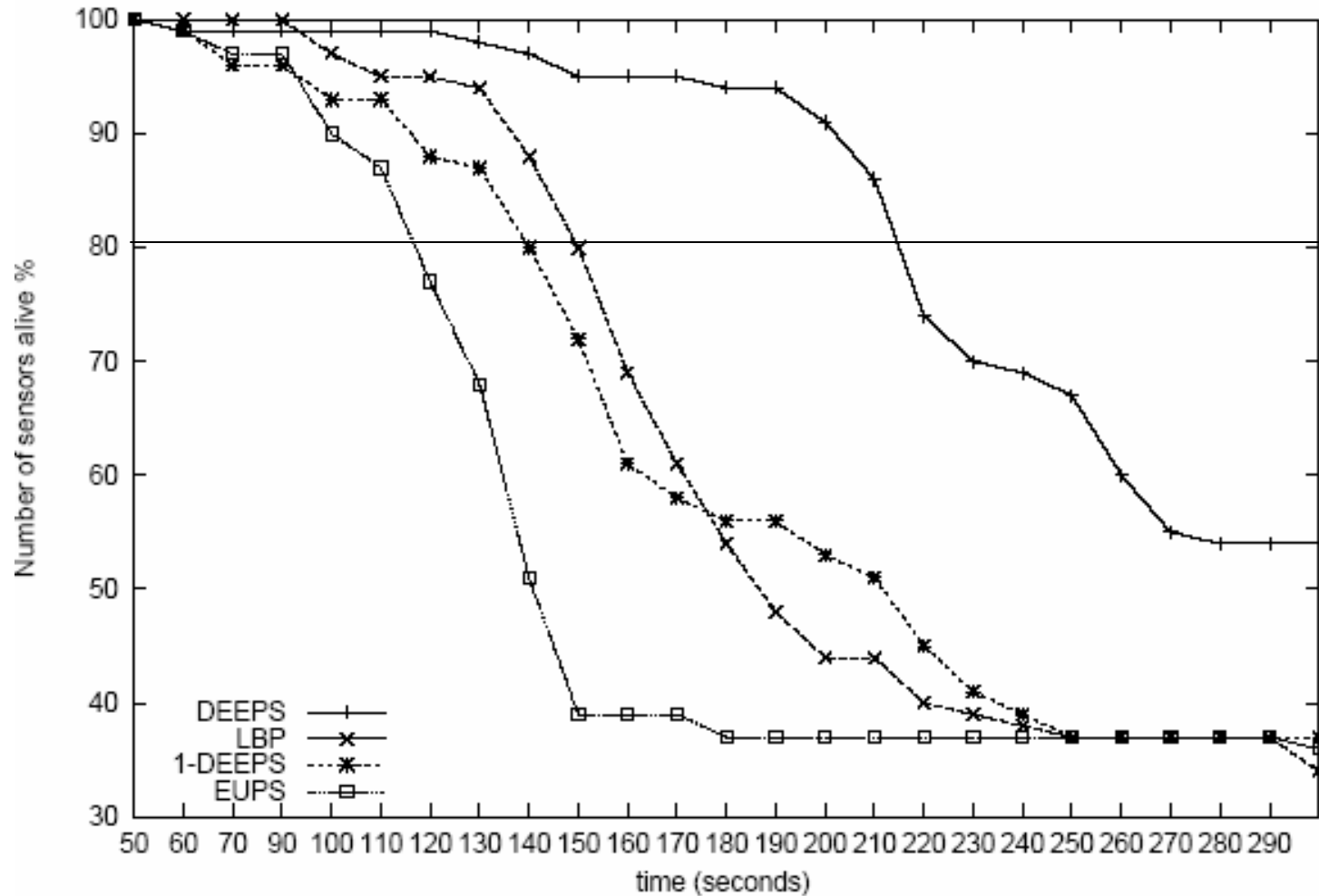
Scenario-1

Covered area for different # of reshuffles



Scenario-1

Number of sensors alive %



Scenario-1

Total energy consumption (J)

